

Network Security (NetSec)

IN2101 – WS 16/17

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Def.: Communications protocol

Problem 1

Problem 2

Recap (Theoretical Comp. Sci.): Chomsky Hierarchy

Problem 3

More on Problem (3): “Weird Machines”

Problem 4

Problem: Mutual Understanding

Examples

Literature and Sources

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Literature and Sources

- Defines the procedure and the format of exchanged messages
- Examples
 - IP
 - TCP
 - UDP
 - HTTP
 - HTTPS
 - SSH
 - ...
- Alice and Bob might speak the same protocol ...
- but do they also have the same understanding?

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`www.paypal.com\0www.zombo.com`

- where `\0` is the C string terminator (NULL character)
- If a browser accidentally uses `strncmp` to validate certificates . . .
- . . . you just got a certificate for `www.paypal.com`

- Alice and Bob spoke the same “protocol”: X.509
- But had a different understanding!
- Alice certified the URL: `www.paypal.com\0www.zombo.com`
- Bob parsed the URL: `www.paypal.com`

- Coder's implicit assumption

Input is well-formed

- Reality

Input is controlled by attacker

- Apply **full** recognition to inputs before processing them!
- Do not scatter recognition throughout your code!



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- My favorite RFC

Content-Length = 1*DIGIT

[...]

Any Content-Length field value greater than or equal to zero is valid. Since there is no predefined limit to the length of a payload, a recipient MUST anticipate potentially large decimal numerals and prevent parsing errors due to integer conversion overflows

- Quiz: Which RFC is this taken from?

- My favorite RFC

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- Quiz: Which RFC is this taken from?
 - 7230, HTTP/1.1 Message Syntax and Routing
- Translation:
 - The length of the content can be arbitrary
 - The length of the Content-Length field can be arbitrary
 - Just parse it right

- What type of grammar is HTTP?
- In the Chomsky hierarchy, at least type 1 – context-sensitive

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- In the Chomsky hierarchy, at least type 1 – context-sensitive
- Are two HTTP parsers equivalent?

UNDECIDABLE

Grammar	Language	Recognized by
Type 3	Regular	Finite state automaton
Type 2	Context-free	Pushdown automaton
Type 1	Context-sensitive	Some weird stuff
Type 0	recursively enumerable	Turing machine

Type 3 \subset Type 2 \subset Type 1 \subset Type 0

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- Remember all those [undecidable](#) problems in theo. comp. sci.?

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- Remember all those [undecidable](#) problems in theo. comp. sci.?
- If the grammar of your protocol is Type 1 or Type 0, you will run into them!

- Don't define Turing-complete protocols
 - Recognizing is undecidable
 - Testing equivalence of different implementations is undecidable
- With Content-Length fields, you easily run into this problem!



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- You are visiting my website

- You are visiting my website
- I host a hidden list of links to the most common porn sites

- You are visiting my website
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- Your browser renders
 - Not visited: blue
 - Visited: purple

- You are visiting my website
- I host a hidden list of links to the most common porn sites
- Your browser renders
 - Not visited: blue
 - Visited: purple
- Using JavaScript, the color of the links is send back to me

- Reduce computing power
- Power that is not there cannot be exploited
- In particular in input handling code



- Complex protocols require complex parsers
- Complex parsers (anything beyond Type 2 and 3) expose almost unlimited computational power to the attacker
- Which leads to “weird machines”

- A weird machine is a machine programmable by an attacker
- Which was not intended or expected by the programmer

- Make your protocol context-free or regular
- And use an appropriate parser
 - Parser generators, parser combinators, . . .
 - `import re` is not an acceptable solution



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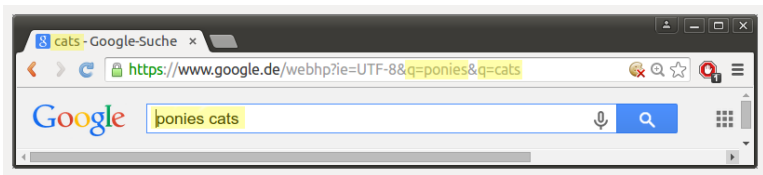
`https://www.google.de/webhp?ie=UTF-8&q=ponies&q=cats`

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- Alice: “*The user asked for ponies*”
- Bob: “*The user asked for cats*”

`https://www.google.de/webhp?ie=UTF-8&q=ponies&q=cats`

- Alice: *“The user asked for ponies”*
- Bob: *“The user asked for cats”*
- Google: *“Let’s go for both (cats preferred)”*



- Entities may have a different understanding of the meaning of a protocol
- In the example
 - Alice recognized the first q parameter
 - Bob recognized the last q parameter

- Messages must be interpreted the same by all participants
- Parsers must be equivalent
- Only decidable for regular and context-free languages



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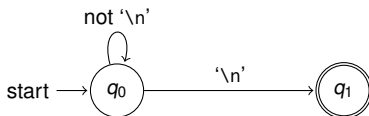
Problem 3

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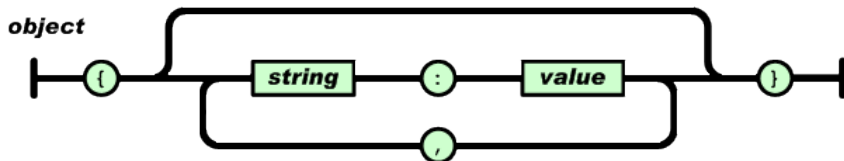
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Literature and Sources

- Familiar from exercises
- Every message is delimited by a '\n'
- Nice library support: `sf.readline()`
- Language is **Regular** (Type 3)



- Context Free (Type 2)



- But: If unique keys are required → no longer context-free

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- Len Sassaman, Meredith L. Patterson, Sergey Bratus, Michael E. Locasto, Anna Shubina, [Security Applications of Formal Language Theory](#), 2013, <http://langsec.org/papers/langsec-tr.pdf>
- <http://langsec.org/>
- Photoshopped protest signs by Kythera of Anevern (www.anevern.com)